I/O PROCESSOR
SERVICE MANUAI

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APPENDIX B

1/O PROCESSOR

## B. 1 GENERAL

The Input/Output Processor (IOP channel (Figure B-1) manages data transfers directly between memory and up to eight multiplexed control-unit/device channels. The IOP contains a pair of address/length control-word registers for each of its eight channels. At the beginning of a transfer for one device, the CPU program uses two WER instructions to load this register pair with the starting address and the block length (Figure B-2). The IOP logic then provides all GP Bus timing ignals to control the data transfers directly between the memory and the CU.
B. 2 For input or output data transfers, a CU signals that it is ready with a Break Request (BR) to the IOP. The IOP makes a Bus Request to obtain control of the GP Bus. The IOP then sends a simulated INR or OTR command to the CU to initiate one word transfer. The INR/OTR command is simulated in that it is generated by the IOP and is not a CPU programmed instruction. The IOP logic updates its control-word'register for each data word. When the block length is counted down to One, the IOP sends End of Record (EOR) to the CU along with the last simulated INR/OTR data transfer. The data block transfer is ended with an SST command and status transfer between the CU and CPU.

## B. 3 The IOP is provided in two versions:

- version with card marked $M X$ is for P852 only.
- version with card marked IOP is for $\mathrm{P} 852 / 856 / 857$

The function and the logic is identical for both versions. The two cards differ in layout and in the location of components. Two sets of logic diagrams are hus provided to show the component locations for each version.


B. $4 \quad \mathrm{CU} /$ Device Priority

Priority of the eight CU/Device channels is established by the connection of the Break Request lines. Refer to paragraph B. 18.

## B. 5 WER Instruction

Two Write External Register (WER) instructions (Figure B-2) are used to load the two control words for a device into a pair of IOP registers. Bits 04, 08-15 of the instructions are sent to the IOP on the MAD lines; bit-15 of each instruction specifies WER-1 or WER-2. The R3 field of each instruction specifies a CPU accumulator (A1-A7) that contains the block-length or start-address control word which is sent to the IOP on the BIO lines.

## B. 6 RER Instruction

A single Read External Register (RER) instruction (Figure B-2) is used when the CPU wants to test the remaining block length of an incomplete data-transfer operation. Bits 04, 08-15 of the instruction are sent to the IOP on the MAD lines; bit $15=0$ agcesses the control-word-1 (block length) register in the IOP. The R3 field of the instruction specifies the CPU accumulator (A1-A7) where the control-word information placed on the BIO lines is to be loaded (bits 4-15).

## B. $7 \quad$ Logic Description and Diagrams

The IOP logic is described in the sequence of its operation, in the section IOP Operating Modes. Operation of some of the more complex logic units (sequensor, ALU, scratchpad) is given also in the section Functional Units. Logic diagrams (Figure B-9, sheets a-e) are provided at the end of the logic description. These diagrams are referenced on the block diagram and in the text by the sheet number, for example: "logic c" refers to Figure B-9, sheet c.
B. 8 IOP OPERATING MODES

The IOP operates in three modes:

- Sampling Mode - - The IOP sequensor logic monitors the IOP status and tests for CU Break Requests or CPU commands.
- CPU Mode - Set upon receipt of a CPU command (WER or RER instruction).

Control-word register information is transferred between the CPU and IOP, with the CPU as master to control the GP BUS.

- Exchange Mode -- Set upon receipt of a Break Request from a CU (with Bus Obtained. One data word is transferred directly between the CU and memory, with the IOP as master to control the GP Bus.


## B. 9 Sampling Mode

The IOP waits in Sampling Mode whenever there is no instruction or data-exchange operation being performed. The sequensor logic is set to the Scan cycle (Figure B-3); timing signals $A P$ and $T l$ are repeated continuously. The AP signal is used to test every 200 ns for instructions from the CPU and Break Requests from the Control Units. Detection of the address-recognized signal AK indicates a CPU command, arid the IOP sets CPU mode with signal NCPU. Detection of a Break Request sets the IOP to exchange mode. The ENB flip-flop (set at the end of an Exchange operation) remains set during Sampling Mode (logic e) to enable a Bus request if a $B R$ is detected.

(BFYBS) AP. ENB
(NCPUZ1) AP. AK. NCPU
(NCW1Z1) AP. MSL.BSYL. TMP

$$
\begin{aligned}
\text { asynchr : } & (B U S R Z 1) B R \\
& (M S L Z 1) \text { BUSR } \cdot \text { OK1 } \\
& (B S Y L Z 1) M S L \cdot \overline{B S Y} \cdot \overline{T R M} \cdot \overline{T P M}
\end{aligned}
$$

Figure B-3 Scan-Cycle Timing
B. 10 L.U Mode

The CPU mode is used by the IOP logic to perform the WER or RER instruction from the CPU. The GPU sends part of the instruction word (including the IOP address) to the IOP (Figure B-2), with timing signal TMEN. The CPU mode is set when the IOP, in Sampling Mode, detects its address-recognized signal AK.
B. 11 Addressing. A WER/RER instruction can address up to 256 external registers with MAD 08-15 (Figure B-2). With bit $08=0$ for 1OP operations, up to 128 external registers can be specified, with 16 registers for each IOP. Bits 12-14 address the register pair for a specific CU/Device channel; bit 15 indicates register-1 or register-2 for the selected device, corresponding to the first or second control word for the channel.
B. 12 CPU-Cycle Operation. CPU-cycle timing is shown in Figure B-4. The IOP compares the address on MAD08-11 with its own address code set by U-links (logic d). An address compare (AKC) is set into the AKCF flip-flop on the rising edge of OSC. The AK signal sets flip-flop NCPU on the rising edge of $A P$ (logic d), if NCPU is not already set. This constraint avoids repeating a CPU cycle. NCPU and $A K$ are used by the Sequensor CPU-cycle: $A P-T 1-B P-T 2-T 3$.
B. 13 For an RER instruction (MADO4 = 1), AK activares BIOVALN (logic d) to gate the control-word register contents onto the BIO lines to the CPU (Scratchpad, logic b). The active RCWN signal inhibits writing into the scratchpad.
B. 14 For a WER instruction (MAD04: 0) , BIOVALN and RCWN are blocked (logic d) and WCWIN is conditioned. The Arithmetic unit (ALU, logic b) is set to logic operation $\vec{B}$ by the selection signals $C W 2,1=1,0$. The control word on the BIO lines is thus connected through the ALU B-input to the scratchpad. The control word is clocked into the addressed scratchpad register by the BP pulse during the CPU-cycle timing.


Note: Times shown in ns
B. 15 The scratchpad-address multiplexer (SPAO-3N) is switched by NCPU so that the MAD lines $12-15$ select the register address. In the middle of T2 time, TRMX is set and the timing-response signal TRMN is sent back to the CPU (logic e)
B. 16 The IOP waits with sequensor-cycle T3 until TMEN from the CPU is terminated. AK is reset on the next OSC after TMEN drops. The loss of AK activates APJN (logic a) which resets TRMX, dropping TRMN, and enables AP to be set on the next OSC. The CPU Mode is finished when AP is set and the IOP is switched to the Sampling Mode.

## B. 17 Exchange Mode

The Exchange Made is used by the IOP logic to perform a data transfer between a CU and memory. The exchange can be either input (CU to memory) or output (memory to CU). The IOP operates as System Master to obtain control of the GP Bus and control the operation. The operation is performed in two logic sequensor cycles: CW1 and CW2. The Exchange Mode is set when the IOP, in Sampling Mode, detects a Break Request (BR) from one of the CUs (Figure B-5).
B. 18 Break Requests. The CU priority is established when the Break Request (BR) lines from the $C U$ s are connected to the IOP, with BROO the highest priority and BRO7 the lowest. The BR lines are examined whenever Enable flip-flop ENB is set (logic e). ENB is reset at the start of an Exchange operation ( $T 1$ of CWI, Figure $B-6$ ) to prevent a higher-priority $B R$ from altering the conditions after an operation has started. ENB is set near the end of the Exchange, and remains set during the Sampling Mode, to enable a $B R$ to initiate a Bus-Request sequence.
B. 19 A Bus Request is initiated if any $B R$ is active when $E N B$ is set. This may happen any time during Sampling Mode, or in the middle of the previous Exchange cycle CW2. Starting a new Bus Request before the previous Exchange has been completed enables data-exchange operations to be linked together to save lime.


Figure B-5 Bus Request Sequence

B. 20 With ENB set, any active $B R$ signal is set into the $F B R$ register (logic e) by the leading edge of AP, which occurs every 200ns during Sampling Mode, or at the end of Exchange cycle CW2. The priority encoder circuit (74148) indicates the code of the highest-priority BR stored in FBR. This code (BRENCO-2) is used to select the external-register (scratchpad) address.

## B. 21 Bus Request Logic. The Bus Control sequence, with the 10 P as

 Master, is shown in Figure $B-5$. BUSRQZIN from the detected $B R$ is gated into the Bus-Request logic (logic e) as BUSRQ if no other operation is active (BUSRQN high). BUSRQ sends BUSRN to request the GP Bus.B. 22 The active-low SPYC response (scan priority chain) sets the OKA flipflop. OKI is received if no higher-priority unit takes control of the Bus. OKI sets the MSL (master selected local) flip-flop, while OKAN blocks the sending of OKO to the next unit on the Bus. MSL sends the Master-Selected signal MSN onto the Bus, and terminates BUSRN. The CPU responds to the end of BUSRN by dropping SPYC and OKI.
B. 23 The IOP may have to wait, with MSL set, until the Bus becomes free of any current operation (TRMN, TPMN, and BSYN all inactive). When the Bus is free, MSL sets BSYLX which puts BSYN onto the Bus to take Bus control for the IOP. The IOP switches to the CWI cycle at the first AP after BSYLX is set.
B. 24 Register Addressing. Two control-word registers must be accessed in the scratchpad during the Exchange Mode. The registers contain data transfer information (control and address) and must be updated during the operation. The scratchpad-address multiplexer (logic b) is switched by NCPU so that the BR piiority encoder and NCW 1 select the register address. The three most-significant bits (SPA1-3) are selected by BRENC2-ON; the least-significant bit (SPAO) is selected by NCW1 (1 for cycle CW1 and 0 for cycle CW2).
B. 25 First Control Word. The first control-word sequence (Figure $B-6$ ) begins with the first AP pulse following the selling of BSYLX. The sequensor

CWI cycle is $A P-T 1-B P-T 2-T 3$, with $A P$ of the CW2 cycle following directly after T3. The NCWI flip-flop (logic d) is set by the leading edge of AP after BSYLX is set.
B. 26 With NCWI set, the ALU mode selected is A minus 1 (for bits 04-15 only). The first control word is gated from the scratchpad into the buffer register (BUF, logic c) by the leading edge of BP. The BUF contents are immediately applied to the ALU where the block-length is decremented and bits 00-03 are transferred directly. If the block length is decremented to One by this operation, ENDN is active. At the end of BP, the updated control word is gated back into the scratchpad.
B. 27 The Input/Output control bit (ALUOIN) and the ENDN signal are set into a two-bit register (logic b) as BUFIN and BUFEOR. These two bits are written into the auxiliary scratchpad as SPINN and SPEORN, and are used to minimize delay in the generation of the MAD signals to the $C U$.
B. 28 The IOP command to the CU is placed on the MAD lines (Figure B-7) via the MADSL multiplexer (logic c). The Input Output control (SPIN) and EOR (SPEOR) bits are sent on MAD04 and 03 respectively. The IOP address (MXADO-2) placed on MAD10-12 is taken directly from the IOP address U-links. The CU address from the priority encoder (BRENCO-2) is sent on MAD13-15.
B. 29 The IOP generates Bus timing signal TMPN (logic e) at the end of time T2 (the first OSC after flip-flop TMPX is set). TMPN goes active 160 ns after the MAD lines are set. The CU uses TMPN to validate the address and contral information on the MAD lines. The CWl cycle for the first control word is completed at the end of T3, before the CU has responded to TMPN.
B. 30 Second Control Word. The sind control-word sequence (tigure B-6) begins with the $A P$ pulse immediately following $T 3$ of $C W 1$. The sequensor $C W 2$ cycle is $A P-T 1^{*}-B P-T 2-T 3^{*}-T 4-T 5-T 6$, with a waiting loop before BP and before T4. Flip-flop NCW2 (logic d) is conditioned by the TMPN signal (TMPXA) and set on the leading edge of AP.

|  |  | SPEOR |  |  | SPIN |  | $\begin{gathered} M X A D \\ 0-2 \\ \vdots \\ 10 \end{gathered}$ | BRENC 0-2 |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| MAD | 128 | 64 | 0 | 3 | 4 | 8 |  | 13 | 15 |
| (IOP - CU) |  |  |  | E | 1 | 00 | IOP |  |  |
|  |  |  |  | R | O |  | Device Address |  |  |

BUF 128,64,00-15


Figure B-7 Bus Content During Exchange Mode (CW1/CW2)
B.31 The second control word is available at the scratchpad output (logic b) as soon as NCW1 goes off. With the control flip-flops NCW1-NCW2 at 0-1, the ALU mode selected is A plus 1 for incrementing the effective character address. To realize A plus 2 for updating the effective word address, CW2WN forces a carry-in to bit 15, resulting in the least-significant bit set to 0 for the evencharacter word addressing.
B. 32 The IOP waits for the CU response TPMN, with the sequensor blocked at $T 1$ before the generation of $B P$. The $C U$ receives TMPN and its address and command on the MAD lines (sent at the end of CW1), places data on the BIO lines (for an input transfer), drops its BR for that exchange, and sends TPMN
back to the IOP. TPMN is clocked into TPMNRA (logic a) on the leading edge of OSC preceeding BP to ensure the guard time of BIO for input transfers.

- B. 33 Before the IOP starts the memory transfer, it ensures that any preceeding memory response signal TRIVN is inactive (important in the crise of linked exchanges). The BP pulse gates the 18-bit effective address (control word 2) from the scratchpad into the buffer (logic c). This memory address is switched via the MADSL multiplexer onto the MAD lines. The selection signal MADSEL is set at BP of CW2. At the same time, the read/write and word/character control signals are sent to the memory via Bus signals WRITE and CHA. These two control bits were read from the first control word and stared as BUFIN (logic b) and BUFCH (logic c) during CWI.
B. 34 The IOP sets TMRX (logice) in the middle of time $T 2$ to send TMRN to the memory. The memory uses TMRN to validate the addiess on MAD (from the IOP) and the data on BIO (from the CU, if input transfer). The IOP must now wait for memory response TRMN.
B. 35 The Bus control flip-flops OKA and MSL (logic e) are reset as a result of TMRN dropping MSN and freeing the Bus for a new selection. The reset occurs when TMRX sets the MSLRF flip-flop, and generates MSLRN. The reset MSLN flip-flop in turn resets MSLRF. BSYN is held active until time $T 6$ so that the IOP maintains Bus control until its operation is complete. Enable flip-flop ENB (logic e) is set on the first OSC after T2 to permit acceptance of a new $B R$ from any of the CUs.
B. 36 If any $B R$ is active when $E N B$ is set, a new Bus Request sequence is initiated (paragraph B-20). If OKI is received without being blocked by another unit on the Bus, another IOP Exchange cycle (CW1-CW2) is linked to the end of this exchange. MSL is again set by OKI and holds BYSN on at time T6 instead of allowing it to be reset as in a normal ending.
B. 37 The IOP waits for the memory response TRMN, with the sequensor blocked at $T 3$ before the generation of T4. The TMPN signal to the CU remains active during the memory-transfer operation to enable the BIO-line data at the CU while TMRN enables BIO at the memory.
B. 38 The TRMN response from memory is set into flip-flop TRMNRA (logic a) on any leading or trailing edge of OSC. Time T4 is then set on the next leading edge of OSC to continue the CW2 cycle. At time T4, flip-flops TMPX and TMPENB (logic e) are reset and TMPN to the CU drops (providing suitable guard time for the CU to read data on the BIO lines). The CU drops TPMN after TMPN goes inactive.
B. 39 For read operations (output, memory to $C U$ ): the memory validates the data on the BIO lines with TRMN, the IOP validates the BIO data at the CU with the trailing edge of TMPN. For write operations (input, CU to memory): the CU validates the BIO-line data with TPMN, the IOP validates the BIO data at the momory with the leading edge of TMRN.


## B. 40 TUNCIIONALUNITS

A block diagram of the IOP is given in Figure B-1. The control words are loaded into the Control Word Registers (scratchpod), via the ALU, by WER instructions.
The CPU reads a Control Word from the scratchpad by means of an RER instruction. During data-transfer operations (Exchange Mode), the Control Words are accessed and updated via the processing loop: scratchpad--buffer--ALU--scratchpad. The address-generation logic sends the CU addiess/command and then the memory address for each data transfer.

## B. 11 Sequensor

The sequensor (logic a) is driven from a constantly-running oscillator to provide the IOP timing sigucals AP, BP, T1, T2, T3, T4, T5, T6. The oscillator frequency at QUARTZ is 20 MHz and OSC is 10 MHz . Sequensor cycles for the different Operoting Modes are:

| Operating Mode | Seq. Cycle |  |
| :---: | :---: | :---: |
| Sampling | Scan | $A P-T 1$, repeated until CPU or Exchange Mode is set. |
| CPU | CPU | AP-T1-BP-T2-T3*, waiting at T3 until the first OSC after TMEN is received from CPU. |
| Exchange | CWI | $A P-T 1-B P-T 2-T 3$, followed by the CW2 cycle. |
| Exchange | CW2 | $A P-T 1^{*}-3 P-T 2-T 3^{*}-T 4-T 5-T 6$, with wating loop before $B P$ and $T 4$. $A P$ of next cycle (Scan $(C W 1)$ follows T4, overlapping T5-T6. |

Signals AP, T1-T4, Th are clocked on the rising edge of OSC. Signals BP and T5 are clocked on the falling edge of OSC and thus shifted one-half OSC cycte from the other sequensor signals. All sequensor timing signals have a duration of one or more 100 ns OSC cycles, as follows

- AP, T2, T4, T5, T6 are always 100 ns .
- BP is 100 ns (CPU cycle) or 200 ns (CW1/CW2 cycles).
- T1 or T3 may last for one or more multiples of 100 ns during the waiting conditions for $B P$ or $T 4$ during the CW 2 cycle.
B. 42 The conditions for the Sequensor timing signals are:
- AP (first pulse of every cycle)

| Condition |  |
| :--- | :--- |
| Scan T1 | Start of Scan or CWI cycle, depending on conditions for AP. <br> CPU T3. AKN <br> Start of Scan cycle; Sequensor vocits of T3 of CPU cycle until <br> AKN indicates the end of the CPU commend, to avoid branching <br> back into a CPU cycle if TMFN is delayed. <br> CW2 T4 |
| Start of CW2 cycle alvays follows T3 of CWI. <br> Start of Scan or new CWI cycle follows $T 4$, although T5-Th of <br> CW2 are not complete. |  |

- T1

| $A P$ | TI follows AP for all cycles |
| :--- | :--- |
| $\overline{B P}$ | Wait at Tl until BP is set (CW2 cycle) |

- BP

| Cycle | Selection |  |  | $\begin{gathered} \text { Set } \\ (B P J) \end{gathered}$ | $\begin{aligned} & \text { Reset } \\ & (\text { BPK }) \end{aligned}$ |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | $\begin{gathered} \text { SO } \\ \text { NCWI } \end{gathered}$ | $\begin{gathered} \mathrm{SI} \\ \mathrm{NCPU} \end{gathered}$ |  |  |  |  |
| CW2 | 0 | 0 | 10 | TPMNRA | 13 | Wait at TI until TPMN received |
| CWI | 1 | 0 | 11 | 1 | T3 |  |
| CPU | 1 | 1 | 12 | 1 | $T 2$ |  |

BPJ enabled by BPJE2N: SCANN.TI.(CW2 TRMN.TMPENB)

- T2 follows BP (CPU, CWI, CW2 cycles)
- T3 follows T2 and stays on until AP (CPU, CWI cycles) or T4 (CW2 cycle) is set.
- T4 follows T3 directly (CPU, CWI cycles) or waits for receipt of TRMN (CW2 cycle).
- T5 (CW2 cycle) ends the command to the CU.
- T6 (CW2 cycle) ensures data guard time for ending the memory transfer.


## B. 43 Arithmetic Unit (ALU)

The ALU (logic b) is used to load control words from the CPU and to update the control words during Exchange Mode. The type 9341 or 74181 ALU circuits operate in one of three modes, controlled by flip-flops NCW1-NCW2, as follows:


* Decrement bits 04-15 only; carry-in to bit 03 blocked by NCW2 so that bits 03-00 are transferred without change.
- CPU mode loads the WER control words via ALU operand B from the BIO lines to the scratchpad.
- Exchange-Mode/CW1-cycle decrements the block-length control word in the processing loop, via ALU operand $A$.
- Exchange-Mode/CW2-cycle increments the effective-address control word in the processing loop, via ALU operand A.
B. 44 Scratchpad

The 16 -word scratchpad (logic b) stores the two control-words for each of the eight Control Units (Figure B-8).
B. 45 Scratchpad Addressing. A type 74157 multiplexer circuit provides - the scratchpad-addressing signals SPAO-3N, as follows:

| Operating <br> Mode | S-Input | Address Source for SPAO-3N |
| :--- | :--- | :--- |
| CPU | NCPU | MADI5-12RF from CPU (MAD lines); aclive low. |
| Exch-CWI | NCPU | NCWI, BRENC2-ON; control-word-1 specified by the <br> CU priority-level code. |
| Exch-CW2 | NCPU | NCWI, BRENC2-ON; control-word-2 specified by the <br> CU priority-level code. |

B. 45 128-64, EOR Bits. A type 74157 selector (SPIN) is used to transfer some bits from one control word to the other. The SPIN selector is controlled by signal SPAON, as follows:

| Cycle | Generation of SPAON, |
| :--- | :--- |
| CPU | MADI5 $=0$, MAD15RF-high, SPAON-high: first control word. |
| CPU | MAD15 = 1, MADI5RF-low, SPAON-low: second control word. |
| CWI | NCWI-high, SPAON-high: first control word. |
| CW2 | NCWI-low, SPAON-low: second control word. |


| SPIN <br> Selection | Source |  |  |
| :---: | :---: | :---: | :---: |
|  | IC | 1 B | IA |
| SPAON-high | ENDN | ALU03N | ALU02N |
| SPAON-Iow | ALU02N | ALU64N | ALUI28N |
| Output: | SPIN02 | SPIN64 | SPINI28 |

- During the first WER instruction (SPAON-high), address bits 128-64 are loaded via SPIN into scratchpad bits 128-64 for the second control word.
- During Exchange-Mode/CWI-cycle (SPAON-high), the decremented block length is tested for a count of ONE, and the.resulting CNDN bit is loaded via SPIN into scratchpad bit 02.
- During operations with the second control word (SPAON-low), the effective address is loaded from BIO (via ALU $\bar{B}$ ) or incremented through the processing loop (via ALU A + 1).



## B. 47 In/Out, EOR bits. The Input/Output and End-of-Record (EOR)

 control bits are sent to the $C U$ on the MAD lines during Exchange Mode, cycle CW1 (Figure B-7). These two bits are obtained via the ALU during the processing loop of the CWI cycle: the In/Out bit from ALUOIN; the EOR bit from testing block-length bits 04-15 for One (ENDN). ALUOIN and ENDN are gated via the type 7475 circuit (as BUFIN and BUFEOR) and loaded into auxiliary scratchpad as SPINN and SPEORN. These two command bits are gated onto MAD04 and 03 along with the control-unit address code during the CWI cycle (which is held until the TPMN response during the CW2 cycle).
## B. 48 Other Logic Units

The remaining logic units are described with the IOP Operating Modes. The following list is a guide to the uses of the different logic units.

| Logic |  | Text Paragraph |  |
| :---: | :---: | :---: | :---: |
| Address Recognition | d | CPU-Cycle | B. 12 |
| CU Request logic | e | Exch/Break Requests | B. 18 |
| Bus Request logic. | e | Exch/Bus Request | B. 21 |
| Bus Request logic. | e | Exch/2nd CW | B. 35 |
| Execution Cycle logic | d | Exch/1st CW | B. 25 |
|  |  | Exch/2nd CW | B. 30 |
| Bus Control logic | e | Exch/1st CW | B. 29 |
|  |  | Exch/2nd CW | B. 34 |
| Buffer | c | Exch/2nd CW | B. 32 |

## B. 49 PHYSICAL

The Units are contained on printed-circuit cards (Figure B-10) of the standard P852/856/857 system size. The parts lists are provided in Table B-1.

The IOP interfaces the other system elements (CPU, control units, and memory) via the GP Bus, via connector-3.

## B. $50 \quad$ U-Links

Adjustable U-links are provided for decoding the IOP address (logic d). The locations of the $U$-links are shown on Figure B-10.
B. 51 Break-Request ( $B R$ ) Signals

The $B R$ signal connections are shown in figure $B-11$. The internal $B R$ signals, from CUs in the same chassis with the IOP, are attached to connector-4B. The external $B R$ signals, from CUs in the extension chassis, are attached to connector 5, and are strapped from connector-4A to 4-B. Priority of the control units is determined by the order of the connections to connector-4B.







Figure B-9a Clock/Sequensor (IOP)




Brackles



Table B-1a Multipley Card Parts List

| Reteience | Description | 12 NC Code |
| :---: | :---: | :---: |
| ```\[ \therefore \quad \therefore .52 .53 .55 \] \[ 0^{\circ}=- \] \[ \text { E5, 20, 57, } 11 . \] \[ \because \] \[ \because \] \[ 2 z=3.03 . \mathrm{c} .5 \mathrm{~s} . \] \[ x: \]``` ```\[ =3 \therefore=1, G 1, G 2,-1, \times 2,11, \mathrm{Ml}, \mathrm{NI} \]``` ```E". \[ \vdots=2 \] \[ =5.23,05,31 . \] \[ =0,=,=z, \mathrm{~N}= \] \[ =\because, 2 c, i t . \] \[ : 5 \] \[ 05 \] \[ \therefore=,-5,34,30, D s, E 5, E s, F s, G s . \] \[ \vdots \] \[ \because \] : \[ \because \] \[ =25 . \] \[ \leq . \] \[ 23 . \] \[ =40 . \] \[ { }^{24} \] \[ \text { P5-15, 20-23, } 51,52 . \] \[ =1-, 24,31-39,47-50,53-58 . \] \[ \because, 10,28,29,30,42,51,52 \] \[ : 1 E, 25,20,27,41,59,50 \] \[ \text { : } 3 \]``` | Printed circuit <br> Integroted circuil 3101 A <br> Integrated circuit 7400 <br> Integrated circuil 7402 <br> Integrated circuit 7404 <br> Integrated circuit 7408 <br> Integrated circuit 7410 <br> Integrated circuit 743" <br> Integrated circuit 7474 <br> Integrated circuil 7475 <br> Integroted circuit 74148 <br> Integrated circuil 74153 <br> Integrated circuir - 4157 <br> Iniegrarea cirevil -4174 <br> Integroted circuil 741:3 <br> 'nteg:ored circuil 9324 <br> Integroted circuit i4181 <br> Integrated cirsuit 1801 <br> Integroted ci:cuit REC OS13 <br> incegroted circuir : 41H30 <br> Integrated circuir 74H53 <br> Inegrated circui: 74500 <br> Integroted circui: 74504 <br> Integrated circuit 74S10 <br> Integrated circuit 74511 <br> Integrated circuit 74520 <br> Integrated circuit 74540 <br> Integrated circuit 74574 <br> Integrated circuit 74S!12 <br> Ternet | 511110005473 |

Table B-1b 10p portstist

| Pelerence | Description | 12 NC Code |
| :---: | :---: | :---: |
| 1.3.86.077 <br> Fs. <br> A7. <br> AI.CG. <br> $B S$. <br> Ab, Q7. <br> C7,F4,H7,K6. <br> w4. <br> B1. <br> A2. <br> H4, NS, Q4, R4, V4. <br> M4. $\begin{aligned} & \text { C1,01,P4,53,54, U1,V3. } \\ & \text { P6, R5, T5,U5,V5. } \end{aligned}$ <br> $M$ : <br> P3,R3,13, i4, U3, W3. <br> G1,H1, J1,K1, L1,M2,N2,P2,Q2,R2,T1. <br> H2,K2,MI,N1,PI,Q1,RI. <br> D6, M6, 2 . <br> DS.FB,KS,N4. <br> B3,C2,C3,M5. <br> B7, D2, H5, P5 . <br> $k 3$ <br> E1. <br> F7. <br> D4. <br> A4, 84, C5,F2,F3,F5,H3,H6,K4. <br> C 1. <br> 243. <br> 284. <br> R4A. <br> $R 45$. <br> $\therefore 46$. <br> $R 52$. <br> Q 30 . <br> 031. <br> 4 3-7, 11-20.22-25,27,28. <br> R2, $1,5,10,21,25,29,32-42,47-51,53-63$ <br> (44-47). <br> C50.63.562 <br> C 67 <br> C 4 <br> C 18 <br> C8.29, 31-43, 54, 65, 68,70-is. <br> FI <br> IRI <br> YI | Printed circuit <br> Integrated circuit 7400 <br> Integrated circuit 7402 <br> Integrated circuit 7404 <br> Integroted circuil 7408 <br> Integrated circuit 7410 <br> Integrated circuit 7437 <br> Integrated circuit 7474 <br> Inregroted circuit 7475 <br> Integrated circuit 74148 <br> Integraled circuit 74153 <br> Integrated circuit 74157 <br> Integrated circuit 74174 <br> Integrated circuit 74175 <br> Integrated circuit 74181 (9341) <br> Integroted circuit 9324 <br> Integrated circuit 3101A <br> Integroted circuit 1801 <br> Integroted eircuit REC 0613 <br> Integrated circuit 74500 <br> Integrated circuit 74504 <br> Integraled circuil 74510 <br> Integrated circuit 74511 <br> Integrated circuit 74520 <br> Integrated circuit 74530 <br> Integrated circuip 74540 <br> Integrated circuir 74564 <br> Integrated circuit 74574 <br> Integrated circuit 74S112 | 511110006052 |



| $1801$ | $7402$ | $7410$ | $7430$ <br> 7430 $\qquad$ |
| :---: | :---: | :---: | :---: |
|  |  | $7411$ | $7437$ |
| $7400$ | $7408$ | $7420$ <br> $7420 \quad 7440$ <br> BUFFER. <br> POSIIIVE y F ABC | $7440$ |


| $7464$  | 74112 <br> 745112 | 74157 | $74181$ |
| :---: | :---: | :---: | :---: |
| $7474$ | $74148$ | $74174$ <br> 74174 $\qquad$ |  |
| $7475$ <br> 7475 |  | $74175$ <br> 74175 | $9324$ <br> REC 0613 |

