

## 2 PROGRAM TESTING

### 2.1 Introduction

A CREDIT program may be input to the PTS system via cards, cassette, flexible disk or console typewriter. After input the source module is held on disk. All the processors and utilities described in Part 2 of this manual read input from disk and write output to disk.

The diagram in page 2.1.2 illustrates the sequence of processes needed to develop and run an executable program from CREDIT source modules.

Each source module is processed separately by the CREDIT Translator. The Translator produces intermediate object code modules. The instructions in these modules use a byte oriented addressing system. Each module may contain references to:

- Labels in the same module
- Literals, formats, key tables and pictures in the same module and/or in the same segment.
- Labels in other CREDIT modules, in the same segment and/or in other segments.
- Assembler application modules.
- Assembler system routines.

The first three types are satisfied by the CREDIT linker.

The remaining types of reference are satisfied by the Linkage Editor. This processor builds an application load module from the following object modules:

- Object modules from the CREDIT linker.
- Assembler application modules (if referenced).
- Assembler system routines.
- CREDIT interpreter.
- CREDIT debugging program (if requested).

CREDIT code cannot be executed directly. Each instruction must be interpreted by the CREDIT Interpreter. The routines within the Interpreter actually perform the functions specified in the CREDIT code. For this reason the Interpreter is built into the Load module by the Linkage Editor.

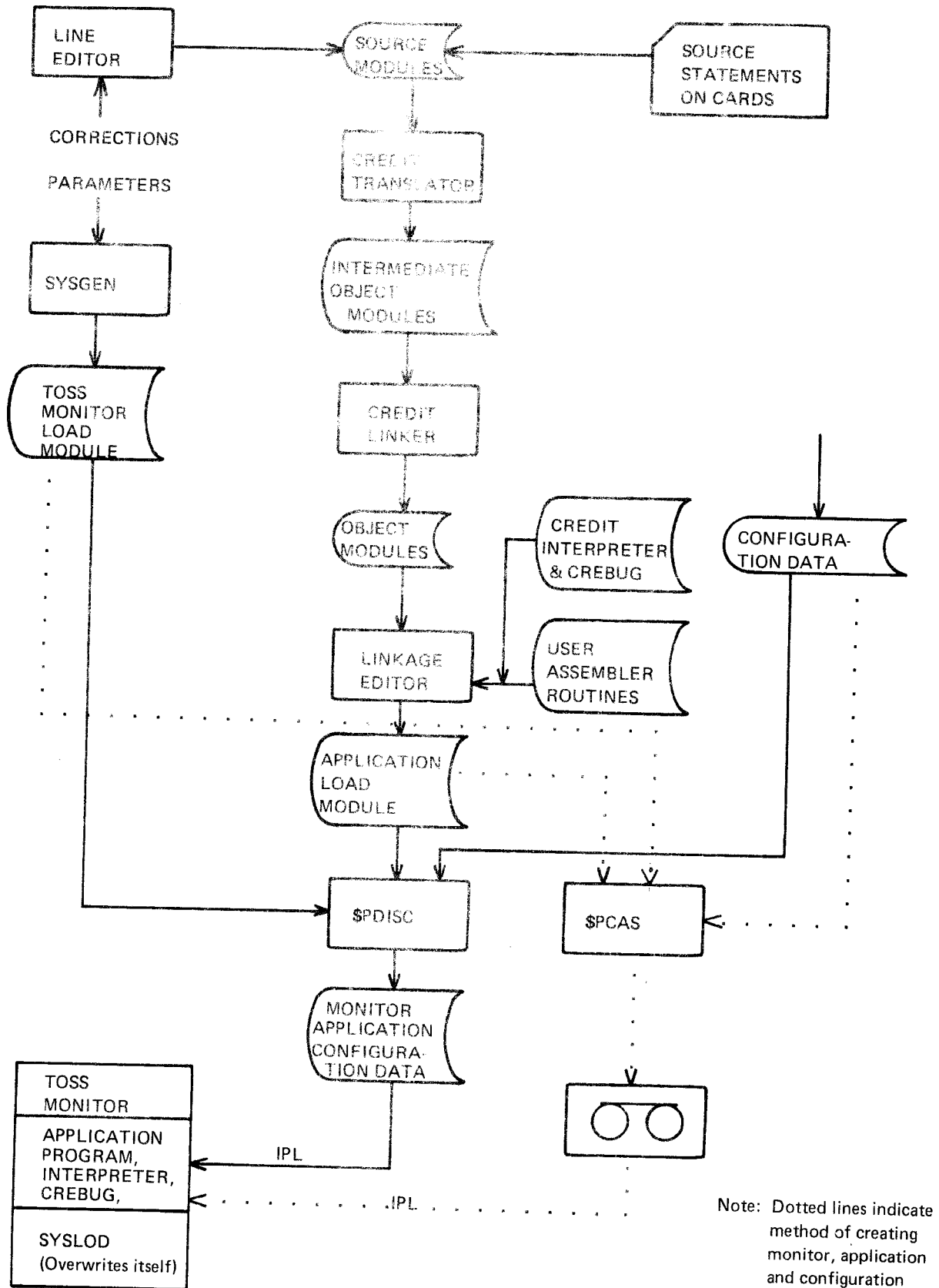
The CREDIT Debugging Program, if required, must also be built into the load module. The Debugging Program is an interactive diagnostic task which, if present, is executed in parallel with the CREDIT program being tested. Via the Debugging Program the programmer can monitor and control the execution of his program.

When the load module is loaded into memory for execution the work blocks, stacks, data set buffers and task control areas required for a particular system, are set up. This is done by the System Loading Program (SYSLOD).

The CREDIT translator, CREDIT linker and linkage editor are all run under the DOS 6800 monitor. The load module produced by the linkage editor, however, must be run under the LPT monitor.

When a particular job, application and application load module, control is handed to the System Loading Program (SYSLOD). When configuration is complete, control is handed to the interpreter. Program testing may then be carried out using one or more work positions. If the CREDIT debugging program is being used, execution can be controlled from the console typewriter.

Development of a CREDIT Application Program



Note: Dotted lines indicate method of creating monitor, application and configuration data for IPL from cassette.

If any application program errors are detected during testing, one or more source modules will have to be corrected. This may be done via the Line Editor – an interactive text editor. Each corrected source module must then be re-processed by the CREDIT Translator. The whole program must then be processed by the CREDIT Linker and Linkage Editor prior to re-testing.

TOSS system software comprises the following components:

- Monitor
- System Loading Program
- CREDIT interpreter
- [CREDIT debugging program]
- [Assembler debugging program]

These software components are not described in a separate manual. Information concerning TOSS System Software which is needed by CREDIT programmers is contained in this manual.

The following software components, though part of DOS6800 System Software, are discussed in this manual:

- CREDIT Translator
- CREDIT Linker

They are discussed in this manual because they are used by CREDIT programmers only. The remaining DOS 6800 System Software components used by CREDIT programmers, notably the Linkage Editor and Line Editor, are described in the DOS6800 System Software PRM (M11).



## 2.2 CREDIT Translator

### 2.2.1 Introduction

The CREDIT Translator is a processor which converts CREDIT source statements into intermediate object code. Source modules are translated separately, resulting in the production of individual object modules. References between object modules and references to external routines etc., are not resolved by the Translator.

Readers of Section 2.2 should be familiar with the following DOS6800 System Software concepts:

- Control Command
- Processor
- EOF mark
- Source input device
- Temporary source file
- Temporary object file

These concepts are explained in the DOS6800 System Software PRM (M11).

### 2.2.2 Running the Translator

Source modules must be read into the System by issuing the control command RDS (read source). RDS will read the source module from the input device (card reader, cassette or console keyboard) and will create a temporary source file. The module must be terminated by an :EOF mark. If the module has been read into the System previously and kept (control command KPF), a RDS command will not be necessary.

It is strongly recommended that all temporary object files are scratched (and kept if necessary) before the Translator is executed. This will ensure that the output object modules will not be corrupted by existing files.

The translator is called into execution by the following control command:

$$\text{TRALI} \left\{ \begin{array}{l} /S \\ \text{name} \end{array} \right\} [,NL]$$

where: /S indicates that the input source module is in the temporary source file.  
name is the name of a source file in the library of the current user identifier. It indicates that the input source module will be found in that file.  
NL indicates that no listing of the module is required. Error messages are always printed.

The intermediate object module created by the Translator is written into the temporary object file. If this file already contains object modules, the following action is taken. If it has not been closed by an EOF mark, the intermediate object module is written *after* the information already held in the file. If it has been closed by an EOF mark, a new temporary object file is created and the old one is deleted.

### 2.2.3 *Translator Listing*

#### 2.2.3.1 *General*

During translation the Translator generates a listing in three parts. Part one contains the CREDIT source statements, intermediate object code and error messages. Parts two and three contain the data item name table and the procedure label table. The following sections describe these parts.

The listing can be suppressed if the NL option is specified on the TRA control command. In this case, only the error messages will be printed.

#### 2.2.3.2 *CREDIT code and Error Messages*

The format of this part is shown in the following example. The example is taken from the procedure division. The data division listing is slightly different. The differences are noted in the explanation which follows.

At the left of the listing under the heading LOC is the location counter. This is a four digit hexadecimal counter which is stepped by one each time a byte of intermediate object code is generated. In the data division a two digit hexadecimal counter called IX (for index) is used.

The next eight items, under the headings OC (operation code) and OPERANDS, comprise the generated interpretive instructions. Each item is a two digit hexadecimal code. The significance of these codes is described for each instruction in the Instruction Reference Section (1.4.8). Object code is not listed in the data division.

The item under the heading LINE is a four digit decimal line counter.

The remaining items are self-explanatory. They comprise CREDIT source statements.

Errors in the source module are reported by the Translator. One of the following messages is printed immediately after the line containing the error:

- 01 Memory overflow (job aborted)
- 02 Sequence error
- 03 Directive missing
- 04 Syntax error
- 05 Length truncation (no error accumulation)
- 06 Multidefined
- 07 Undefined
- 08 Unexpected value
- 09 Undefined type
- 0A Unexpected type
- 0B Illegal constant
- 0C Lit pool overflow
- 0D Label missing
- 0E Illegal value def
- 0F Illegal const length
- 10 Illegal const type
- 11 Too many blocks
- 12 Too many data items.
- 13 Block size overflow
- 14 Too many datasets
- 15 Too many parameters
- 16 Too many start stmts
- 17 Illegal dimension.
- 18 Too many values.
- 19 Out of range.
- 1A Unspecified parameter.
- 1B Parameterlist overflow.

In addition to the error message an asterisk is printed to show the position at which the error occurred.

An error count is maintained by the Translator and is printed after the END directive.

If a fatal error occurs (I/O error, table overflow, etc), the source input is read until an EOF mark is encountered and the following message is printed:

FATAL ERROR HAS OCCURRED. NO OBJECT CODE PRODUCED.

The object file is then deleted.

#### 2.2.3.3 Data Item Name Table

The data item name table is listed immediately after the CREDIT code and error messages. For each data item declared in a work block it contains the following:

NAME	This is the data item identifier.
REF	This is the index number assigned by the Translator. Index numbers are printed to the left of the data item declarations, under the heading 'IX'.
TYPE	This is the data item type specified in the data item declaration. The following mnemonics are printed under the TYPE heading: BCD (decimal), BIN (binary), BOL (boolean) and STR (string). A letter U following one of these mnemonics indicates that the data item is not referenced in this module.

#### 2.2.3.4 Procedure Label Table

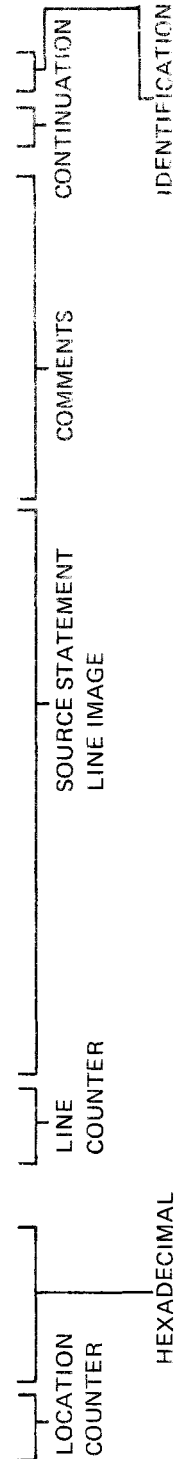
The procedure label table is printed immediately after the data item name table. For each identifier appearing in the procedure division it contains the following:

NAME	This is an identifier specified by the programmer in the procedure division, or it is the name of a System routine referred to by the generated object code.
REF	This may be an index to a format list, a key table or an external table. It may be a value specified in an equate directive. It may be the contents of the location counter (if the identifier belongs to an instruction or PROC directive).
TYPE	This indicates the type of identifier. The following mnemonics are used: ADR Address of an instruction EQU Equate directive EXT External label FOR Format list KEY Key table PRO PROC directive FLB Format label (address of a format item) FTB Format table

A letter U following one of these mnemonics indicates that the identifier is not referred to in this module.

CREDIT TRANSLATOR REL 3.1 780520 \* PROCEDURE DIVISION IDENT OPCLOS \* 060477 \* PAGE 0006

LOC	OC	OPERANDS	LINE LABEL	OPCOD	OPERANDS	COMMENT	C	1D
			0336					
			0337 *		EJECT			
			0338 *					
			0339 *		ENTRY TERMINAL CLOSE MODULE			
			0340 *					
			0341 *					
			0342	OPCL05				
0000	30	XX 85 FF	0343	EDWRT DSTPJT,FRM003		PRINT ASTERIX		
0004	30	XX 85 FF	0344	EDWRT DSTPJT,FRM014		PRINT CLOSE		
			0345					
			0346	OPC010				
0008	XX	KK LL LL	0347	PERF READN,KTAB3,=W'4',=W'4'		READ OPERID		
000C	5F	06	0348	SB OPC010				
			0349					
000E	20	60 42 05	0350	CB EQ,CASHID,INPUT,OPC100		CORRECT IDENT (TY		
0012	XX	12 LL	0351	PERF ERROR,TLAMP3,KTAB4				
0015	5F	0F	0352	SB OPC010				



HEXADECIMAL REPRESENTATION OF THE INSTRUCTION

- RR = Reference in this module
- LL = Literal constant
- XX = External reference or operation code
- FF = Reference to a format list
- KK = Reference to a key table

CREDIT TRANSLATOR LISTING EXAMPLE



## 2.3 CREDIT Memory Management Linker

### 2.3.1 Introduction

The CREDIT memory management linker is a three pass processor which converts intermediate object code, produced by the CREDIT translator into object code which can be processed by the Linkage Editor. The CREDIT linker is capable of linking both unsegmented and segmented programs.

Intermediate object modules may contain references to:

- Labels in the same module.
- Literal constants, formats, key tables, pictures in the same module.
- Labels in other CREDIT modules in the same segment.
- Labels in other CREDIT modules in other segments.
- Assembler application modules.
- Assembler System routines.

The first three types of reference, when present in the same segment, are satisfied by the CREDIT linker.

The remaining types of reference must be satisfied by the Linkage Editor.

To build up the segments, different possibilities exist which are the same when using extended main memory, secondary memory or a combination of both.

Readers should be familiar with the following DOS6800 System Software concepts:

- Control command
- User library
- User identifier
- Temporary object file

These concepts are explained in the DOS6800 System Software PRM (M11).

### 2.3.2 Building up segments

After translation of the different CREDIT modules a number of intermediate object modules have been created. All these modules together building up a CREDIT application, are input to the CREDIT linker (TLK).

The CREDIT linker has to know which modules should be contained in the segments. This is specified by the user by means of the ordering of the modules as input to the linker. In the TLK command is a parameter (n or mK) defining the maximum segment size to be used.

However, the NOD command (node) can be used to force an immediate end of a segment and also to define the segment as main memory resident (NOD  $\rightarrow$  R) or belonging to the common area, segment 00 (NOD  $\rightarrow$  C). The common area, segment 00, is always present in main memory and will contain the data division, the interpreter, assembler sub-routines and /or user routines. The size of the common area is variable and not dependent on the size parameter in the TLK command. Segment 00 is automatically created.

When the NOD command is used without specifying R or C as parameter, the segment will be disk resident. When the NOD command is not used, the segments will be disk resident.

The output from the CREDIT linker is placed in temporary object file (/O) and divided into different modules as: data division, common part and segments. The user can extend the common part, segment 00, with the NOD LJC command. The modules are grouped into segments in the order in which they are included with the INC command. The segments are numbered from 1 upward.

Some examples showing the use of NOD and TLK for different systems.

a. System with 64K Byte main memory.

```
INC  MOD1
INC  MOD2
INC  MOD3
TLK  U,M,X
```

b. System with 64K Byte main memory and use of secondary memory. (Disk, flexible disk).

	Size
INC  MOD1	(3.5Kbytes)
INC  MOD2,  USER1	(0.7Kbytes)
INC  MOD3,  USER2	(0.5Kbytes)
INC  MOD4	(2Kbytes)
INC  MOD5	(2Kbytes)
INC  MOD6,  USER3	(1Kbyte)
INC  MOD7	(2Kbytes)
TLK  U,M,4K	

Four segments are created by the linker, and all are disc resident. (Segment zero always resides in main memory.)

Segment 1	Contains	MOD1	(3.5K)
Segment 2	Contains	MOD2, MOD3, MOD4	(3.2K)
Segment 3	Contains	MOD5, MOD6	(3K)
Segment 4	Contains	MOD7	(2K)

By means of altering the sequence of the INC commands, the user can optimize his program segments. In this example only 50% of segment 4 is filled.

When e.g. MOD5 must be present in segment 0, the following sequence of commands has to be specified:

	Size
NOD  C	
INC  MOD5	(2Kbytes)
NOD	
INC  MOD1	(3.5Kbytes)
INC  MOD2,  USER1	(0.7Kbytes)
INC  MOD3,  USER2	(0.5Kbytes)
INC  MOD4	(2Kbytes)
INC  MOD6,  USER3	(1Kbyte)
INC  MOD7	(2Kbytes)
TLK  U,M,4K	

Segment 1 contains: MOD1  
 Segment 2 contains: MOD2, MOD3, MOD4  
 Segment 3 contains: MOD6, MOD7,.

MOD5 is now included in the common area, segment zero.

When also MOD1 must be main memory resident, but not in segment 0, then the following command sequence can be used:

		Size
NOD	C	(2Kbytes)
INC	MOD5	(2Kbytes)
NOD	R	
INC	MOD1	(3.5Kbytes)
NOD		
INC	MOD2, USER1	(0.7Kbytes)
INC	MOD3, USER2	(0.5Kbytes)
INC	MOD4	(2Kbytes)
INC	MOD6, USER3	(1Kbyte)
INC	MOD7	(2Kbytes)
TLK	U,M,4K	
Segment 1 contains:	MOD1	(3.5Kbytes)
Segment 2 contains:	MOD2,MOD3,MOD4	(3.2Kbytes)
Segment 3 contains:	MOD6, MOD7	(3Kbytes)

c. System with extended main memory, up to 256Kbytes.

		Size
INC	MOD1	(3.5Kbytes)
INC	MOD2, USER1	(0.7Kbytes)
INC	MOD3, USER2	(0.5Kbytes)
INC	MOD4	(2Kbytes)
INC	MOD5	(2Kbytes)
INC	MOD6, USER3	(1Kbyte)
INC	MOD7	(2Kbytes)
TLK	U,M,4K	

Four segments are created by the linker, and are assumed to be disk resident. Because an extended main memory is used, all segments will be main memory resident. The composition of the segments is as mentioned in example b.

The system loader SYSLOD will discover the difference when a system with extended main memory, secondary memory or a combination of both is used.

d. Systems with extended main memory (up to 256 bytes) and secondary memory.

Examples b) and c) may be combined.

### 2.3.3 Running linker

The CREDIT linker reads intermediate object modules from temporary object file and from the library of the current user identifier. The syntax of the TLK command is:

TLK u [N|S|U][,X] [,M] [,n[,mK]

N The system or user /OBJCT files do not need to be scanned.

U Only the user /OBJCT files will be scanned.

S Only the system /OBJCT file has to be scanned.

Default value: Both /OBJCT files will be scanned.

The user /OBJCT file will be scanned first, then the system /OBJCT file and then the user /OBJCT file again.

X Indicates that a cross reference listing is required.

Default value: No cross reference will be printed.

M The listing of the map, which consists of a listing of the module names and the relative start addresses, address pools and statistics per segment.

Default value: No map will be printed.

n The required segment size in bytes.

mK The required segment size in K bytes. (m x 1024 bytes)

Default value: The program will be unsegmented.

### 2.3.3.1 CREDIT modules in the system library

When CREDIT modules have to be linked from the System library (USERID:SYSTEM), first a Generate Object Directory command (GOD) must be executed for this library before the TLK command can be executed. (In any other user than SYSTEM.)

The following listings are produced by the CREDIT memory management linker per segment:

#### — Segment 0 (Common part)

load map

long branch table

call table

perform table

literal pool

key table pool

picture pool

format pool

} can be excluded by not using  
'M' in the TLK command.

Linker statics for this segment

#### — Segment n

loadmap

long branch table

perform table

literal pool

picture pool

format pool

} can be excluded by not using  
'M' in the TLK command.

Linker statics for this segment

#### — Total

segment map

cross reference

} can be excluded by not using  
'X' in the TLK command.

Linker statics for the whole program.

The load map includes a list of error reports. Error reports will be listed even if the load map listing has not been requested.

2.3.3.2 Load Map

The load map indicates the displacement of each module within a segment. It also contains the linker (TLK) error reports. The format of the load map is shown in the following example:

```

-----
* CREDIT CODE LINKER PRR 4.1 790410 * LOAD MAP SEGMENT 02
-----

LOC      MODULE  ERROR      COMMENT

000E     MOD3      TRA 4.1  99-99-99 F1 01111
006D     MODUL4      TRA 4.1  99-99-99 F1 01111
009D     MODUL6      TRA 4.1  99-99-99 F1 01111
00B9     MODUL7      TRA 4.1  99-99-99 F1 01111
00CB     MODUL8      TRA 4.1  99-99-99 F1 01111
    
```

where: LOC is the displacement of the module within the segment.  
 MODULE is the module name.  
 ERROR is the error number followed by a type, number and clear text.

Error type may be:

- E – User Error
- I – Internal error or input inconsistency
- W – Warning, no updating of error counter.

The following error reports may be printed:

ERROR NUMBER	ERROR TYPE	Additional Information	Text (Significance)
0	I		END OF MEMORY No more work space available
1	E		SYMBOL TYPE CONFLICT LB, CALL or PERF mixed up
2	I	XXXX	ILLEGAL INPUT XXXX is a hexadecimal presentation of 1st and 3rd character in cluster. Input from translator not expected.
3	I	XXXX	LOAD ADR INCORRECT XXXX is a hexadecimal presentation of load address from the cluster.
4	W	DDDD	UNREFERENCED LITERAL DDDD is a decimal presentation of the number of unreferenced literals.

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ERROR NUMBER	ERROR TYPE	Additional Information	Text (Significance)
5	I	DDDD	UNDEFINED LITERAL DDDD is a decimal presentation of the number of undefined literals.
6	E		NO START ADDR Start address declaration not found in the data division.
7	E		DBL DEF MODULES Double defined modules.
8	E	DDDD	UNSATISFIED EXTERNAL DDDD is a decimal presentation of the number of unsatisfied externals. The unsatisfied externals are printed on the load map.
9	I	XXXX	MODULE LENGTH ERROR XXXX is a hexadecimal presentation of the difference between requested and available workspace.
10	E	DDDD	TRANSLATION ERROR DDDD is a decimal presentation of the number of translating errors.
11	E	XXXX	WRONG TRANSLATOR RELEASE XXXX is a hexadecimal representation of the lowest acceptable level of the CREDIT translator, in the form RRLL RR = Release number LL = Level number
12	E	C	ADDRESS TABLE OVERFLOW C is a character representing the address type L (Long Branch), C (Call) or P (Perform).
13	E	C	LITERAL DISPLACEMENT OVERFLOW C is a character representing the literal type L (Literal), K (key table), P (picture), or F (format).
14	E	C	TOO MANY LITERALS C is a character representing the literal type L (literal), K (key table), P (picture), or F (format).
15	E		FORMAT LENGTH ERROR
16	E		MULTI DEF ENTRY Entry name defined in more than one module.
17	E	C	NOD TYPE ERROR C is a character representing the NOD type. NOD type not C, D or R.

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ERROR NUMBER	ERROR TYPE	Additional Information	Text (Significance)
18	E	XXXX	MODULE LONGER THAN SEGMENT SIZE XXXX is a hexadecimal representation of the module length. Increase segment size.
19	I		IDENT MISSING
20	E		ADDRESSING MODE CONFLICT One byte or two bytes addressing mode of literals mixed up.
21	E		NOD SEQUENCE ERROR The NON record "NOD C" does not appear in the beginning of the object input.

2.3.3.3 Call table

This table contains all references to external routines (CALL instruction) which could not be satisfied by the TLK command. Each time a reference is encountered in the intermediate code, the linkage editor (LKE command), replaces it by an "index value" which points to the called address in the call table. During execution of the application program, the interpreter refers to the call table for actual destination addresses. The format of the call table is shown in the following example:

-----  
 \* CREDIT CODE LINKER PRR 4.1 790410 \* CALL TABLE SEGMENT 00  
 -----

LOC	DATA	IX	SYMBOL	DEFINED
0002	****	01	T:ASSI	
0004	****	02	T:KI	
0006	****	03	T:EDWR	
0008	****	04	T:DSC1	
000A	****	05	T:NKI	
000C	****	06	T:RREA	
000E	****	07	T:RWRI	

where: LOC is the displacement of each table entry within segment zero.  
 DATA is the called address relative to the start of segment zero. It is generated by the link map editor and is therefore not specified in the listing.  
 IX is the index value for the entry (maximum index is X'FF')  
 SYMBOL is the name of the external routine.  
 DEFINED is printed in the table.

2.3.3.4 Long branch table

In order to reduce the amount of memory required for a long branch instruction, linker (TLK) generates a table of destination addresses. Each time a long branch is encountered in the intermediate code, the linker places the destination address (i.e. segment number and the address to be branched to) in the long branch table.

The three byte destination address in the long branch instruction is replaced by a one byte "index value" which points to the destination address in the long branch table. During execution of the application program, the interpreter refers to the long branch table for actual destination addresses. The format of the long branch table is shown in the following example:

```
-----
* CREDIT CODE LINKER PRR 4.1 790410 * LB TABLE SEGMENT 01
-----
```

LOC	DATA	IX	SYMBOL	DEFINED
0170	01 007A	01		MODULE5
0174	01 0054	02		MODULE5
0178	01 0118	03		MODULE5
017C	01 009D	04		MODULE5
0180	01 00F8	05		MODULE5
0184	01 00E3	06		MODULE5
0188	01 00C0	07		MODULE5
018C	01 010F	08		MODULE5

where: LOC is the displacement of each table entry within the segment.  
 DATA is the destination address of the long branch. The first two digits specify the segment number and the next four specify the displacement within this segment. The difference between the four digit hexadecimal value, and the relevant module start address shown in the load map, gives the address of the destination within that module.



- IX is the index value used in the long branch instructions. It starts at one. (Maximum index is X'FF').
- SYMBOL is the statement identifier of the first instruction (location counter = 0) in the module containing the destination of the branch.
- DEFINED is the name of the module containing the destination.

2.3.3.5 Perform table

This table contains the address of each CREDIT subroutine which is called (PERF instruction) within this segment. It has the same layout as the long branch table. Each time a perform to a CREDIT subroutine is encountered, in the intermediate object code the subroutine name is replaced by an "index value" which points to the subroutine address in the perform table. The format of the perform table is shown in the following example:

-----  
 \* CREDIT CODE LINKER PRR 4.1 790410 \* PERFORM TABLE SEGMENT 01  
 -----

LOC	DATA	IX	SYMBOL	DEFINED
0192	01 0037	01	VDU1	MOD1
0196	XX XXXX	02	K81	
019A	XX XXXX	03	VDU2	
019E	01 005E	04	K82	MODUL5
01A2	XX XXXX	05	DISC	
01A6	01 0030	06	GTP1	MOD2
01AA	XX XXXX	07	VDU3	
01AE	XX XXXX	08	VDU4	
01B2	01 0046	09	GTP2	MOD2

where: LOC is the displacement of each table entry within the segment.  
 DATA is the destination address of the perform. The first two digits specify the segment number and the next four specify the displacement within this segment. The difference between the four digit hexadecimal value, and the relevant module start address shown in the loadmap, gives the address of the destination within that module.

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- IX is the index value. It starts at one (maximum index is X'FF').
- SYMBOL is the name of the sub-routine. It only appears when the sub-routine is not in the same module as the perform instruction.
- DEFINED is the name of the module which contains the sub-routine.

2.3.3.6 *Literal pool*

The literal pool contains all the literals used in this segment. Each time a literal is encountered in the intermediate code it is replaced by an "index value" which points to the literal in the literal pool. The format of the literal pool is shown in the following example:

-----  
 \* CREDIT CODE LINKER PRR 411 290450 \* LITERAL POOL SEGMENT 02  
 -----

IX	TYPE	LOC	DATA
10	BIN	0008	0000
11	BIN	000A	0004
12	BIN	000C	0006
13	BIN	000E	0008
14	BIN	0010	0009
15	BIN	0012	0016
16	BIN	0014	0033
17	BIN	0016	0040
18	BIN	0018	0042
19	BIN	001A	0200
1A	BIN	001C	0410
1B	BIN	001E	1410
1C	STR	001F	07
1D	STR	0021	2028
1E	STR	0023	2030
1F	STR	0025	203107
20	STR	0028	4141411E
21	STR	002C	4141421E
22	STR	0100	4141431E
23	STR	0104	4141441E
24	STR	0108	4141451E

- where: IX is the index value. It starts at 10 or 4100 (maximum index is X'FFF' or X'4FFF').
- TYPE indicates the value type of the literal. The following mnemonics are used:  
 BIN for value types X and W.  
 BCD for value type D.  
 STR for value type C.
- LOC is the displacement of each literal within the segment.
- DATA is the hexadecimal representation of the literal.

### 2.3.3.7 *Picture pool*

The picture pool contains all picture strings used in this segment. Each time a reference to a picture string is encountered in the intermediate code, it is replaced by an "index value" which points to the picture string in the pool. The format of the picture pool is shown in the following example:

```
-----
* CREDIT CODE LINKER PRR 4.1 790410 * PICTURE POOL SEGMENT 01
-----
```

IX	TYPE	LOC	DATA
10	PIC	01D1	393939
11	PIC	01D4	5A5A5A5A5A5A5A392C3939
12	PIC	01DF	3939452D3939452D39393939
13	PIC	01EB	5A5A565A5A5A565A5A392C39392B

where: IX is the index value. It starts at 10 or 5100 (maximum index is X'FF' or X'5FFF').  
 TYPE indicates the entry is a picture string (PIC).  
 LOC is the displacement of each picture string within the segment.  
 DATA is the hexadecimal representation of the picture string.

### 2.3.3.8 *Keytable pool*

The keytable pool contains all keytables used in the application program and is located in segment zero. Each time a reference to a keytable is encountered in the intermediate code, it is replaced by an "index value" which points to the keytable in the pool. The format of the keytable pool is shown in the following example:

```
-----
* CREDIT CODE LINKER PRR 4.1 790410 * KEYTABLE POOL SEGMENT 00
-----
```

IX	TYPE	LOC	DATA
10	KEY	0010	031E1D19

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where: IX is the index value. It starts at 10 or 6100 (maximum index is X'FF' or X'7FFF').  
 TYPE indicates the entry is a keytable. (KEY)  
 LOC is the displacement of the keytable within segment zero.  
 DATA is the hexadecimal representation of the keytable. First character in the keytable is the length indicator.

2.3.3.9 Format pool

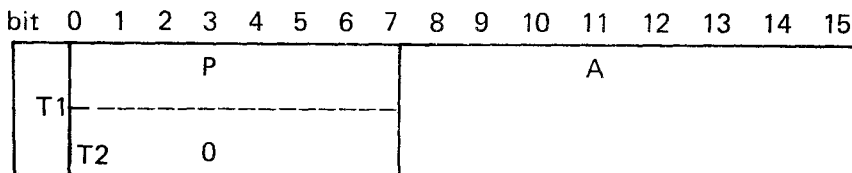
The format pool contains all format lists used in the segment. Each time a reference to a format list is encountered in the intermediate code, it is replaced by an "index value" which points to the format list in the pool. The format of the format pool is shown in the following example:

-----  
 \* CREDIT CODE LINKER PRR 4.1 790410 \* FORMAT POOL SEGMENT 01 \* DATE 2857 \* PAGE 9 \*  
 -----

IX	TYPE	LOC	DATA
10	FMT	01F9	C118C027
11	FMT	01F0	C118C026
12	FMT	0201	C118C025
13	FMT	0205	C1181322
14	FMT	0207	C11CC30F415554484F524954592E2E2E2E2E3A9E20C305444154453A
15	FMT	0225	C11C9420C3192A2A5452414E53414354494F4E2043414E43454C4C45442A2AE8C11CA82AE8C11CA82A
16	FMT	024E	C11CA82AE8C11CC31F454E44204F46204441592C5345525649434520444953434F4E54494E554544E8C11CA82AE8C11CA82A
17	FMT	0280	C11CC30F5452414E53414354494F4E2E2E2E3A10219E20C305444154453A1220E8C11CC11CC30E4F50455241544F5220434F44453AC02A9E20C30D4F4646494345204E4F3A313233E8C11CC3054E414D453AC0258620C3074144524553533AC026E8C11CC305434954593AC0278920C3084143434F554E54204E4F3AC027E8C11CC307414D4F554E543A11229020C30C4E45572042414C414E43453A1123

where: IX is the index value. It starts at 10 or 7100 (maximum index is X'FF' or X'7FFF').  
 TYPE indicates that the entry is a format list (FMT) or format table (FTB). The layout of FMT entries is explained below.  
 LOC is the displacement of each format list in the pool within the segment.  
 DATA is the hexadecimal representation of the format list or format table.

Each word in an FMT entry has the following layout:



Depending on the T1 and T2 bit, fields P and A or 0 and A have the following meaning:

- T1 = 0; P field contains an index to a picture string (FMEL).
- A field refers to decimal-data-item

T1 = 1; 0 field contains a six-bit value, indicating how many times the character in the A-field has to be copied. (FILLR).

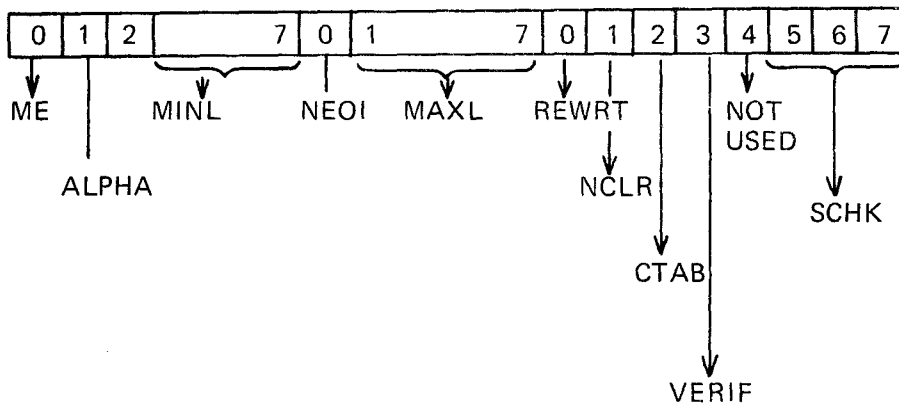
T1 = 1, T2 = 1;

Contents 0-field	Significance
00/01	A-field contains a reference to a string-data-item or literal (FCOPY).
03	A-field and following bytes contain ISO-7 characters (FTEXT).
04	A-field contains a tabulation value. (FTAB).
08	Character X'1F' edited into the buffer. A-field not used (FHIGH).
09	Character X'1E' edited into the buffer. A-field not used (FLOW).
0A	Character X'12' edited into the buffer. A-field not used (FUL).
0B	Character X'13' edited into the buffer. A-field not used (FNUL).
10	A-field contains a reference to a binary or decimal-data-item. The byte following the A-field contains a displacement. (FBZ).
11	A-field contains a reference to a binary or decimal-data-item. The byte following the A-field contains a displacement. (FBP).
12	A-field contains a reference to a binary or decimal-data-item. The byte following the A-field contains a displacement (FBN).
14	A-field contains a reference to a binary or decimal-data-item. The byte following the A-field contains a displacement (FBNZ).
15	A-field contains a reference to a binary or decimal-data-item. The byte following the A-field contains a displacement (FBNP).
16	A-field contains a reference to a binary or decimal-data-item. The byte following the A-field contains a displacement (FBNN).
18	A-field contains a displacement (FB).
1A	A-field contains a reference to a boolean-data-item. The byte following the A-field contains a displacement (FBF).
1B	A-field contains a reference to a boolean-data-item. The byte following the A-field contains a displacement (FBT).
1C	A-field contains a reference to a binary-data-item. (FCW).
1D	A-field contains a reference to a literal (FCW).
1F	A-field contains a reference to a subformat list (FLINK).
20	A-field not used (FSL).
21	A-field not used (FNL).
28	A-field not used (FEOR).
29	A-field not used (FEXIT).

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- 2C A-field contains the tabulation position (FINP).
- 2E A-field contains the tabulation position, the following byte contains the right halfword of the application (APPL) control field (FINP).
- 2F A-field contains the tabulation position, the following 2 bytes contain in sequence:
  - a) left halfword of the application (APPL) control field
  - b) right halfword of the application (APPL) control field (FINP).
- 30 A-field contains the tabulation position. The three following bytes constitute the standard input control field (FKI).

Layout standard control field (FKI).



- 32 A-field contains the tabulation position. The following byte contains the right halfword of the application (APPL) control field. The next three bytes are the standard control field bytes (see also 30) (FKI).
- 33 A-field contains the tabulation position. The following bytes contain in sequence:
  - a) left halfword of the application (APPL) control field
  - b) right halfword of the application (APPL) control field
  - c), d) and e) are the standard control field bytes (see also 30) (FKI).
- 38 A-field contains the tabulation position. The following bytes contain in sequence :
  - a) duplication data-item (DUPL) reference.
  - b), c) and d) are the standard control field bytes (see also 30) (FKI).
- 3A A-field contains the tabulation position. The following bytes contain in sequence:
  - a) right halfword of the application (APPL) control field
  - b) duplication data-item (DUPL) reference
  - c), d) and e) are the standard control field bytes (see also 30) (FKI).

3B

A-field contains the tabulation position the following bytes contain in sequence:

- a) left halfword of the application (APPL) control field
- b) right halfword of the application (APPL) control field
- c) duplication data-item (DUPL) reference
- d), e) and f) are the standard control field bytes (see also 30) (FKI).

2.3.3.10 Linker statistics per segment

The format of the linker statistics listing per segment, is shown in the following example. The contents of the listing are self-explanatory.

```
-----
* CREDIT CODE LINKER PRR 4.1 790410 * LINKER STATISTICS   SEGMENT 00
-----
ALL VALUES DECIMAL

      LB TABLE:      0 BYTES.      0 ENTRIES
      CALL TABLE:   14 BYTES.      7 ENTRIES
      PERFORM TABLE: 0 BYTES.      0 ENTRIES

LITERAL DESCRIPTOR TABLE: 0 BYTES. 0 ENTRIES
PICTURE DESCRIPTOR TABLE: 0 BYTES. 0 ENTRIES
KEYTABLE DESCRIPTOR TABLE: 4 BYTES. 1 ENTRIES
FORMAT DESCRIPTOR TABLE: 0 BYTES. 0 ENTRIES

      LITERAL POOL SIZE: 0 BYTES
      PICTURE POOL SIZE: 0 BYTES
      KEYTABLE POOL SIZE: 4 BYTES
      FORMAT POOL SIZE: 0 BYTES

INTERPRETABLE CODE SIZE: 0 BYTES
PROGRAM LENGTH:        40 BYTES

NUMBER OF ERRORS      0
```

2.3.3.11 Segment map

This map gives a listing of the number of segments, the number of modules contained in a segment and the number of bytes per segment. The format of the segment map is shown in the following example:

```
-----
* CREDIT CODE LINKER PRR 4.1 790410 * SEGMENT MAP
-----

      S E G M E N T
NUMBER TYPE LENGTH  USAGE  N U M B E R  O F
                                M O D U L E S  E R R O R S

00      C           40
01      D          914   91 %           4           0
```

where: NUMBER is the segment number.

TYPE

indicates:

C = common part (segment zero)

R = main memory resident

D = disk resident

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LENGTH number of bytes contained in this segment (program length).  
 USAGE a filling percentage of the segment, related to the size option in the TLK command.  
 MODULES number of modules contained in the segment.  
 ERRORS number of errors per segment.

2.3.3.12 Address cross reference listing

This listing provides a cross reference between statement/subroutine identifiers in the procedure division and the modules/segments in which they are referenced. The format of the address cross reference listing is shown in the following example:

-----  
 \* CREDIT CODE LINKER PRR 4.1 790410 \* CROSS REFERENCE LISTING \* DATE 2857 \* PAGE  
 -----

SYMBOL	TYPE	VALUE	SEG-DEFINED	REFERENCES
DISC	P	02 0070	02-MODUL6	01-MAIN (1)
GO	B S	01 000E	01-MAIN	02-MODUL6 (1) 02-MODUL7 (1) 02-MODUL8 (1)
GTP1	P	01 0030	01-MOD2	01-MAIN (1)
GTP2	P	01 0046	01-MOD2	01-MAIN (1)
GTP3	P	01 004C	01-MOD2	01-MAIN (1)
KB1	P	02 000E	02-MOD3	01-MAIN (1)
KB2	P	01 0052	01-MODUL5	01-MAIN (1)
T:ASS1	C			01-MAIN (1)
T:DSC1	C			01-MODUL5 (8) 02-MOD3 (4) 02-MODUL4 (3) 02-MODUL6 (1)
T:EDWR	C			02-MODUL7 (1) 01-MOD1 (1) 01-MOD2 (3) 01-MODUL5 (4) 02-MOD3 (1)
T:KI	C			02-MODUL4 (1) 02-MODUL6 (1) 02-MODUL7 (1) 02-MODUL8 (1)
T:NKI	C			01-MAIN (1) 01-MODUL5 (3) 02-MOD3 (2)
T:RREA	C			01-MODUL5 (2)
T:RWRI	C			01-MODUL5 (1) 02-MODUL6 (1)
V0U1	P	01 0037	01-MOD1	01-MAIN (1)
V0U2	P	02 0060	02-MODUL4	01-MAIN (1)
V0U3	P	02 00B9	02-MODUL7	01-MAIN (1)
V0U4	P	02 00C8	02-MODUL8	01-MAIN (1)

where: SYMBOL is the statement/subroutine identifier in the procedure division.  
 TYPE indicates type of instruction in which "symbol" is used.  
 C = CALL  
 P = Perform  
 B = Branch  
 S = Start point.  
 VALUE displacement of "symbol" in the referenced segment. The first two digits specify the segment number and the next four specify the displacement within this segment.



SEG-DEFINED segment number and module name which contains "symbol".  
REFERENCES are the segment numbers and module names, containing  
references to "symbol". The number of references in each  
module appears in brackets after the module name.

2.3.3.13 *Linker statistics total*

The format of the linker statistics total is shown in the following example. The contents of the listing are self-explanatory.

-----  
\* CREDIT CODE LINKER PRR 4.1 790410 \* LINKER STATISTICS TOTAL  
-----

ALL VALUES DECIMAL

INTERPRETABLE CODE SIZE: 547 BYTES  
PROGRAM LENGTH: 1568 BYTES  
  
AVAILABLE WORKSPACE: 23874 BYTES  
USED WORKSPACE: 2934 BYTES, 12 %  
UNUSED WORKSPACE: 20940 BYTES  
MAX WORKSPACE PER MODULE: 370 BYTES  
  
NUMBER OF ERRORS 0  
  
PROG ELAPSED TIME: 00H-02M-01S-520MS-

